Amendments to the Claims

This listing of claims will replace all prior version, and listings, of claims in the application:

Listing of Claims:

1. (Original) A method by which a first computing entity having an RSA key pair (N_A,e_A), (N_A,d_A) digitally signs and encrypts a message data string, m, for decryption by a second computing entity having an RSA key pair (N_B,e_B), (N_B,d_B), where $|N_A| = |N_B| = k$ and $m \in \{0,1\}^n$, and $k = n + k_0 + k_1$ for integers k_0 and k_1, the method comprising:

a) selecting an integer $r \in \{0,1\}^{k_0}$,

b) computing:

$w \leftarrow H(C_1(\text{at least } m \text{ and } r))$

where $H : \{0,1\}^{n+k_0} \rightarrow \{0,1\}^{k_1}$, and $C_1()$ is a deterministic combination function,

c) computing:

$s \leftarrow Enc(w, C_2(\text{at least } m \text{ and } r))$

where $Enc()$ is a symmetric-key encryption function using w as key, and $C_2()$ is a reversible combination function;

steps a) to c) being repeated as necessary to obtain $s||\omega \leq N_A$; and then

d) signing by computing:

$c' \leftarrow (C_3(\text{at least } s \text{ and } w))^{d_A} \bmod N_A$

where $C_3()$ is a reversible combination function; and

e) if $c' \leq N_B$, encrypting c' by computing:

$$c = c'^{e_B} \bmod N_B.$$

2. (Original) A method according to claim 1, wherein if $c' > N_B$ following step d), the most significant bit of c' is removed to obtain a new c' which is then encrypted by computing:

$$c = c'^{e_B} \bmod N_B.$$

3. (Original) A method according to claim 1, wherein if $c' > N_B$ following step d), steps a) to d) are repeated as necessary to obtain $c' \leq N_B$ whereupon c' is encrypted by computing:

$$c = c'^{c_B} \bmod N_B$$

4. (Original) A method according to claim 1, wherein r is selected at random.

5. (Original) A method according to claim 1, wherein the function $C_1()$ is a concatenation function.

6. (Original) A method according to claim 1, wherein the function $C_2()$ is a concatenation function.

7. (Original) A method according to claim 1, wherein the function $C_3()$ is a concatenation function.

8. (Original) A method according to claim 1, wherein the functions $C_1()$, $C_2()$, $C_3()$ are all concatenation functions.

9. (Previously presented) A method according to claim 1, wherein the symmetric-key encryption function *Enc()* effects at least the following operations:

- forming a hash of the key w;
- forming an exclusive-OR of the hash of w with the output of the combination function $C_2()$.

10. (Previously presented) Apparatus for carrying out the method of claim 1.

11. (Previously presented) A computer-readable medium storing a computer program arranged to condition a program-controlled computer, when executed by the latter, to carry out the method of claim 1.

12. (Previously presented) A method according to claim 1, wherein the second computing entity on receiving c:

(f) computes:

$$c' \leftarrow c^{d_B} \bmod N_B$$

and, provided $c' \leq N_A$, proceeds to the next step;

(g) computes:

$$c'^{e_A} \bmod N_A$$

with the result being subject to a reverse of the combination function $C_3()$ whereby to recover at least: s and w;

(h) computes:

$$Dec(w, s)$$

where $Dec()$ is a symmetric-key decryption function complimenting $Enc()$, with the result being subject to a reverse of the combination function $C_2()$ whereby to recover at least: m and r;

(i) checks that the message m is from the first computing entity by checking that:

$$w = H(C_1(\text{at least } m \text{ and } r)) .$$

13. (Previously presented) A system comprising a first computing entity, a second computing entity, and a communications network for communicating the first and second entities, the system being arranged to implement the method of claim 12.

14. (Previously presented) A method according to claim 2, wherein the second computing entity on receiving c:

(f) computes:

$$c' \leftarrow c^{d_B} \bmod N_B,$$

and, provided $c' \leq N_A$, proceeds to the next step;

(g) computes:

$$c'^{e_A} \bmod N_A$$

with the result being subject to a reverse of the combination function $C_3()$ whereby to recover at least: s and w;

(h) computes,

$$Dec(w, s)$$

where $Dec()$ is a symmetric-key decryption function complimenting $Enc()$, with the result being subject to a reverse of the combination function $C_2()$ whereby to recover at least: m and r;

(i) checks that the message m is from the first computing entity by checking that:

$$w = H(C_1(\text{at least } m \text{ and } r));$$

j) where the check carried out in step (i) fails, computes a new value for c' as:

$$c' \leftarrow c' + 2^{k-l}$$

and, provided $c' \leq N_A$, repeats once steps (g) to (i).

15. (Previously presented) A system comprising a first computing entity, a second computing entity, and a communications network for communicating the first and second entities, the system being arranged to implement the method of claim 14.

16. (Previously presented) A method by which a second computing entity having an RSA key pair (N_B, e_B), (N_B, d_B), decrypts and authenticates a ciphertext c that is purportedly a signed and encrypted form produced by a first computing entity of a message data string m, the first computing entity having an RSA key pair (N_A,e_A), (N_A,d_A) where $|N_A| = |N_B| = k$ and $m \in \{0,1\}^n$, and $k = n + k_0 + k_1$ for integers k_0 and k_1; the second computing entity on receiving c:

(a) computes:

$$c' \leftarrow c^{d_B} \bmod N_B$$

and proceeds to the next step provided that $c' \leq N_A$;

(b) computes:

$$c'^{e_A} \bmod N_A$$

with at least quantities *s* and *w* being recovered from the result;

(c) computes:

Dec(w,s)

where *Dec()* is a symmetric-key decryption function complimenting *Enc()*, with at least quantities *m* and *r* being recovered from the result;

(d) checks that the message *m* is from the first computing entity by checking that:

$w = H(C_1$(at least *m* and *r*))

where $H : \{0,1\}^{n+k_0} \rightarrow \{0,1\}^{k_1}$ and $C_1()$ is a deterministic combination function.

17. (Previously presented) A method according to claim 16, wherein the function $C_1()$ is a concatenation function.

18. (Previously presented) A method according to claim 16, wherein the symmetric-key decryption function *Dec()* effects at least the followings operations:

- forming a hash of the key *w*;
- forming an exclusive-OR of the hash of *w* with *s*.

19. (Previously presented) Apparatus for carrying out the method of claim 16.

20. (Previously presented) A computer-readable medium storing a computer program arranged to condition a program-controlled. computer, when executed by the latter, to carry out the method of claim 16.

21. (Currently amended) A method by which a first computing entity having an RSA key pair (N_A,e_A), (N_A,d_A) digitally signs and encrypts a message data string, m, for decryption by a second computing entity having an RSA key pair (N_B, e_B), (N_B, d_B), where $|N_A| = |N_B| = k$ and $m \in \{0,1\}^n$, and $k = n + k_0 + k_1$ for integers k_0 and k_1 <u>even</u>, the method comprising:

a) selecting an integer $r \in \{0,1\}^{k_0}$,

b) forming the hash $\omega = H(m \,||\, r)$ where $H : \{0,1\}^{n+k_0} \rightarrow \{0,1\}^{k_1}$, and

c) forming the hash $s = G(\omega) \oplus (m \,||\, r)$ where $G : \{0,1\}^{k_1} \rightarrow \{0,1\}^{n+k_0}$;

steps a) to c) being repeated as necessary to obtain $s \,||\, \omega \leq N_A$; and then

d) signing by forming $c' = (s \,||\, \omega)^{d_A} \bmod N_A$; and, if $c' > N_B$,

removing the most significant bit of c' to obtain a new c'; and then

e) encrypting c' by forming $c = c'^{e_B} \bmod N_B$.

22. (Previously presented) The method as claimed in claim 21 in which r is selected at random.

23. (Previously presented) A computer storage medium having stored thereon a computer program readable by a general-purpose computer, the computer program including instructions for said general purpose computer to configure it for implementing the steps of the method of claim 21.

24. (Currently amended) A method by which a first computing entity having an RSA key pair (N_A,e_A), (N_A,d_A) digitally signs and encrypts a message data string, m, for decryption by a second computing entity having an RSA key pair (N_B,e_B), (N_B,d_B) where $|N_A| = |N_B| = k$ and $m \in \{0,1\}^n$, and $k = n + k_0 + k_1$ for integers k_0 and k_1 <u>even</u>; the method comprising:

a) selecting an integer $r \in \{0,1\}^{k_0}$,

b) forming the hash $\omega = H(m \,||\, r)$ where $H : \{0,1\}^{n+k_0} \rightarrow \{0,1\}^{k_1}$, and

c) forming the hash $s = G(\omega) \oplus (m \,||\, r)$ where $G : \{0,1\}^{k_1} \rightarrow \{0,1\}^{n+k_0}$;

steps a) to c) being repeated as necessary to obtain $s \,||\, \omega \leq N_A$ and then

steps a) to c) being repeated as necessary to obtain $s \| \omega \leq N_A$ and then

d) signing by forming $c' = (s \| \omega)^{d_A} \bmod N_A$;

steps a0 to d) being repeated as necessary to obtain $c' < N_B$, and then

e) encrypting c' by forming $c = c'^{e_B} \bmod N_B$.

25. (Previously presented) The method as claimed in claim 24 in which r is selected at random.

26. (Previously presented) A computer storage medium having stored thereon a computer program readable by a general-purpose computer, the computer program including instructions for said general purpose computer to configure it for implementing the steps of the method of claim 24.